

Development of Methods for Research of Electric Power System Flexibility

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Abstract. This paper presents deterministic methods developed to study the flexibility of an electric power system. They rely on the proposition that an electric power system is flexible if power balance is maintained at a considered time. These methods are aimed at determining the combination of the largest loads, which, when exceeded a little, disturb power balance at studied nodes. The paper presents two methods: brute-force optimization and nonlinear optimization. Results obtained using the first method are taken to be a reference for verification of nonlinear optimization output.

1 Introduction

In terms of control of an electric power system (EPS), the flexibility of the EPS that has generating equipment with specific maneuverability characteristics is closely related to its ability to maintain frequency and voltage in the system under uncertainty and variability [1]. Thermal and hydroelectric power plants, which can quickly ramp up and ramp down the load, provide the flexibility of the EPS on the generation side. A variety of load management techniques that have emerged owing to the development of new technologies solve the flexibility problem on the demand side. With the adoption of wind and solar farms, energy storage is becoming an important tool for ensuring flexibility.

A prerequisite for ensuring the EPS flexibility is operating reserves available in the system. The considered reserves or sources of flexibility are:

1. Operating reserves [2], [3], [4].
2. Demand management [5].
3. Energy storage systems [6].

Reserve is used in the case of an unplanned increase or decrease in load. A large number of sources of variable generation (wind, solar) in the EPS require the placement of upward and downward reserves [2]. Authors of [3] provide an overview of the operating reserves used in the USA and Europe. In [4], a methodology for determining the minimum required reserves of Russia's EPSs is given.

Demand-side management is applied to adjust residential load [5], service sector load [6], and industrial load [7].

Energy storage devices are used to store and deliver power during a certain time. Energy storage technologies are based on different physical principles. The following classification of energy storage devices is given in [8]:

- Mechanical: flywheels, hydraulic accumulators, pneumatic accumulators.
- Electric: capacitors and supercapacitors.
- Electrochemical: storage batteries, hydrogen fuel cells, nano-ion cells.

Author of [9] describes Superconducting Magnetic Energy Storage (SMES), which stores energy in a magnetic field created by direct current in a coil with zero electrical resistance, cooled below a characteristic critical temperature.

Researchers in many countries are studying the issues of the flexibility margin, presence, and absence in the power system. There are currently probabilistic and deterministic methods for determining flexibility.

In [11], a deterministic method was proposed to determine the largest variation range of uncertainties at which the power system remains flexible for a specified time within acceptable cost. The flexibility metric is calculated by comparing the obtained range with the target range. In [10], the flexibility residual, which is the difference between the available flexibility and the expected load ramps is calculated for each observation and horizon. Then, the probability that the residual flexibility will be less than zero is determined, which means the probability of insufficient resources in the system. In [11], the flexibility of thermostatically controlled loads (TCLs) when integrated into system-level operation and control is calculated. The authors propose a geometric approach to modeling the aggregate flexibility of TCLs. The set of valid power profiles of individual TCLs is represented by a polyhedron. Aggregated flexibility is calculated as the Minkowski sum. The authors developed an optimization algorithm for approximating polynomial by homotheties of a given convex set represented by a virtual battery model.

The insufficient ramping resource expectation (IRRE) metric to estimate flexibility is calculated in [12]. For each direction and time horizon, a probability distribution of IRRE is formed.

This paper presents deterministic methods developed to study the EPS flexibility. The structure of the paper is as follows. The second section discusses modeling the flexibility of EPS facilities and modeling the archive of loads. The third section presents the ideas of methods, the objective function of calculating the maximum loads and constraints. The fourth section provides a detailed description of the methods for calculating flexibility. The fifth section presents the research results. The sixth section is the conclusion.

2 Modeling the flexibility of EPS elements and load archive

Model of generator flexibility of a conventional station

The flexibility available from each generator is determined by the power that can be generated over the considered time horizon and is calculated by the formula [12]

$$F_g = V_{i+} * (t - (1-b) * S_i), \quad (1)$$

where V_{i+} is load ramp time (MW/min), t is the considered time horizon, S_i is the start-up time (hour), b is the binary on-line variable when a generator is on $b=1$.

Model of battery flexibility

The flexibility available from the battery is determined by its state of charge. If the battery is charged within the specified limits

$$SOC_{\min} < SOC(t) < SOC_{\max}, \quad (2)$$

then the power output is calculated by the formula:

$$F_B = P_{max}, \quad (3)$$

otherwise:

$$F_B = 0. \quad (4)$$

Model of system flexibility

System flexibility is determined as total flexibility available from all units of flexibility

$$F_S = \sum_1^m F_g + \sum_1^n F_B, \quad (5)$$

where m is the number of generators at conventional stations, n is the number of batteries.

Modeling of load archive

The load at each given node i is calculated by the formula [13]:

$$P_i(z_i) = z_i P_i^{min} + (1 - z_i) P_i^{max}, \quad (6)$$

where $0 \leq z_i \leq 1$, P_i^{max} is the upper limit of load at node i , P_i^{min} is the lower limit of load at node i .

The archive of loads is formed according to the following algorithm:

1. Set the minimum and maximum values of the active load. Form vectors P^{min} and P^{max} , $P^{min} = (P_1^{min}, P_2^{min} \dots P_l^{min} \dots P_R^{min})$, $P^{max} = (P_1^{max}, P_2^{max} \dots P_l^{max} \dots P_R^{max})$, where R is the number of given load nodes.

2. Set vector $z = (z_1, z_2, \dots, z_i, \dots, z_R)$. Specify the number of steps N , which determines the size of the archive. Calculate the step of changing the load by the formula

$$step = 1 / N. \quad (7)$$

Initial condition: $z = (0)$ is the zero vector, $k=1$ is the step number.

3. Calculate the value of load by (6).

4. If $k = N * R$, go to item 7, otherwise $k = k + 1$, go to item 5.

5. Calculate z^{k+1} by

$$z^{k+1} = z^k + step. \quad (8)$$

6. Go to item 3.

7. Determine all possible load values.

8. The end. The result is a modeled archive of loads P^{LOAD} , the dimension of the archive is $[L \times R]$ where

$$L = C_N^R.$$

3 The idea of the methods. Objective function and constraints

This paper presents deterministic methods based on the proposition that an EPS is flexible if a power balance is maintained at the considered time. An increase in the load leads to a decrease in the flexibility of the system, this is why one of the key points in the analysis of the EPS flexibility is the availability of information about the maximum possible loads. The developed methods are aimed at determining the combination of maximum loads which, when exceeded a little, disturb the power balance at the studied nodes.

The objective function is the maximum sum of the differences between the predicted and simulated loads at nodes with uncertainty over a given time. It is written as follows:

$$\sum_{i=1}^r (\bar{P}_i - P_i(z_i)) = \sum_{i=1}^r \Delta P_i(z_i) \rightarrow max \quad (9)$$

where r is the number of nodes with uncertainty.

For clarity of presentation of the constraints used to solve this problem, all nodes are divided into three types:

- Uncontrolled nodes. Generator nodes where control actions are not performed or load nodes at which there is no uncertainty P^{CONST} ;
- Controlled nodes. Generator nodes where the control actions P^{CA} are performed;
- Nodes with uncertainty. Load nodes at which power changes.

The constraints are as follows:

$$\Delta P_j = 0, \quad (10)$$

$$P_{i-j} < P_{i-j}^{max}, \quad (11)$$

$$P_i^{min} < P_i^{CA} < P_i^{max}, \quad (12)$$

$$0 \leq z_i \leq 1, \quad (13)$$

where in (9) \bar{P}_i is the forecast (pseudo measurement) of active power at node i , which has uncertainty; $P_i(z_i)$ is the relationship between active power and value z , which is responsible for a change in the value of power at node i . Constraint (10) is the power balance at node j (any type of node), or the power balance at EPS, (11) is the constraint on line transfer capability; P_{i-j}^{max} is the capability limit of transmission line $i-j$, (12) limits the range of control actions at the controlled node, (13) is a constraint on the parameters of optimization.

4 A detailed description of the developed methods

The paper presents two methods for determining flexibility:

1. Brute-force optimization.
2. Nonlinear optimization.

4.1 The brute-force optimization

The brute force optimization is used to process all combinations of possible loads in EPS to determine load combinations that, when slightly exceeded, make the system inflexible.

The brute force optimization algorithm is described below.

1. Start. The vector of injections is $P = (P^{CONST}, P^{CA}, P^L)$. Calculate a load flow solution (LFS). $P^{ref} = P^L$, where P^L is a load at the nodes with uncertainty at a given time. Initial conditions are $P^{rab} = P^{ref}$; $i = 1$.

2. Perform the control action P^{CA} according to $P^{LOAD}(i)$

3. Form the vector of injection $P = (P^{CONST}, P^{CA}, P^{LOAD}(i))$.

4. Calculate a load flow solution.

5. If the process has converged, go to the next step. Otherwise, go to step 9.

6. Check the constraints (formulas 10-13).

7. If the constraints have been satisfied, go to the next step. Otherwise, go to step 9.

8. Compare the vectors $(P^{rab} - P^{ref}) < (P^{LOAD}(i) - P^{ref})$. When the condition is

met, save the vector $P^{LOAD}(i)$, $P^{rab} = P^{LOAD}(i)$. Use Euclidean distance and distance of Chebyshev to compare the two vectors.

9. If $i = i + 1$. $i = L$, go to step 10, otherwise, go to step 2.

10. The end. The result: $P^{LOAD}(i)$ is a combination of the largest loads in EPS, which are possible under the given conditions.

4.2 Nonlinear optimization

Nonlinear minimization refers to the problem of nonlinear programming and is performed in Matlab. As a result, the values of optimization parameters used for the calculation of active power P_i^{calc} at the nodes with uncertainty are determined. In this study, the optimization parameter is z (formula (6)). Therefore, the objective function (9) and constraints (10), (11) should be written using the parameter z . Constraint (12) is taken into account by the objective function.

The objective function.

Each element of (9), taking into account (6), can be written as:

$$\Delta P_i(z_i) = P_i - P_i(z_i) = \bar{P}_i - P_i^{max} + z_i(P_i^{max} - P_i^{min}) = D_i + F_i z_i \quad (14)$$

$$\bar{P}_i - P_i^{max} = D_i, \quad (15)$$

$$P_i^{max} - P_i^{min} = F_i. \quad (16)$$

D_i, F_i remain the constant values during the optimization process.

The objective function can be written as follows

$$D_i + z_i F_i + \dots + D_R + z_R F_R + z_r F_r \rightarrow \max \quad (17)$$

and after excluding all constant values it has a compact form:

$$\sum_{i=1}^R (-F_i z_i) - F_A z_A \rightarrow \max. \quad (18)$$

Constraints

After some transformation, power balance in EPS

$$\sum_{i=1}^{un} P_i + \sum_{j=1}^{n-un} P_j(z_j) = 0 \quad (19)$$

can be written as follows

$$\sum_1^{n-un} F_j z_j = \sum_1^{u-un} P_j^{max} + \sum_1^{un} P_i, \quad (20)$$

where n is the number of nodes in EPS, un is the number of uncontrolled nodes.

Building the balance and transmission constraints, which are needed to ensure that all state variables be within their limits, requires the values of power flows in the lines. In this study, power flows in the lines are calculated using the PTDF (power transfer distribution factor) method [15]. PTDFs describe how active power flows in lines are changed if power injection at the node is increased or decreased.

The power transfer distribution factor in the line limited by nodes i, j is calculated in advance as follows

$$k_{i-j} = \Delta P_{i-j} / \Delta P_A(z_A), \quad (21)$$

$$\Delta P_A(z_A) = \sum_{i=1}^R \Delta P_i(z_i), \quad (22)$$

where ΔP_A is an increase (decrease) of the active power at the node, where the control action is performed ΔP_{i-j} is an increase (decrease) of the active power flow in line $i-j$, ΔP_i is an increase (decrease) of active power at the node with uncertainty.

For the problem of nonlinear optimization, the coefficients k_{i-j} are the initial data.

The power balance at nodes with uncertainty is compiled as a balance of power increments

$$\Delta P_i + \sum_{j=1}^G \Delta P_{i-j} = b, \quad (23)$$

$$\Delta P_{i-j} = \Delta P_A(z_A) k_{i-j}. \quad (24)$$

where ΔP_A is an increase (decrease) of active power at node A , ΔP_{i-j} is power flow increments, G is the number of adjacent nodes, b is convergence tolerance. Given (14), (24), the constraint at node i is written as:

$$F_A Z_A \sum_{j=1}^G k_{i-j} + F_i Z_i + (P_i - P_i^{max}) \sum_{j=1}^G k_{i-j} = b. \quad (25)$$

Active power flows in lines are monitored according to (11)

$$P_{i-j} + k_{i-j} \Delta P_A(z_A) < P_{i-j}^{max} \quad (26)$$

when transferring constant values to the right-hand side (taking into account (14) for $i = A$), constraint (26) has the form

$$F_A k_{i-j} z_A < P_{i-j}^{max} - P_{i-j} - k_{i-j} (P_A - P_A^{max}). \quad (27)$$

4.3 Determination of EPS flexibility

The flexibility of EPS is determined as follows:

$$F_S = \sum_{i=1}^R (P_i^{calc} - P_i^{forec}). \quad (28)$$

where P_i^{calc} is a calculated (simulated) value of the active load at node i ; P_i^{forec} is a forecast of active load at node i . If

$$F_S > 0,$$

then EPS is flexible.

5 Case study

5.1 Describing a test scheme and scenario

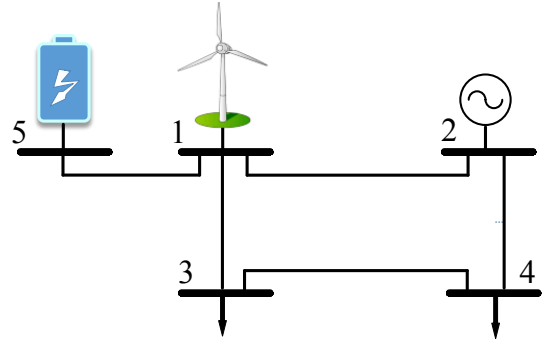


Fig.1. Test scheme

Calculations are performed using a scheme consisting of 5 nodes and 5 lines (figure 1). Nodes 3 and 4 are the nodes with uncertainty. Node 1 stands for a wind farm. Node 5 is a battery. Node 2 (conventional plant) is a controlled node where control actions are power generation required to ensure balance in the EPS, given the forecast of generation from the wind farm and the power supplied by the battery. Nodes 1 and 5 are considered to be uncontrolled.

The calculations are performed according to the scenario: it is necessary to determine a combination of the largest loads at nodes 3, 4 four minutes ahead with known:

- forecasts of load values at nodes 3 and 4 (P^{forec}), which are assumed to be the lower load limits;
- forecast of active power output at the wind farm;
- forecast of active power output at the battery;
- maximum load values, which are the upper load limits;
- maximum value of active power generation at node 2;
- capacity limits of transmission lines.

It is assumed that it takes 4 minutes for the entire available reserve at the conventional plant to be switched on and that the battery produces maximum power.

Table 1. Initial data (MW).

| Number of nodes | p^{max} | p^{min} | p^{forec} |
|-----------------|-----------|-----------|-------------|
| 1 | | | |
| 2 | 32 | 20 | 20 |
| 3 | 37 | 23 | 23 |
| 4 | 23 | 13 | 13 |
| 5 | | | |

The flexibility of a 5-node EPS is calculated by two methods: the brute-force optimization and non-linear optimization. The results of the first method are taken to be a reference.

5.2 The brute-force optimization applying

Using this method, the vector of active loads is determined among 900 pre-created vectors that differs as much as possible from the forecasted loads when the following constraints are met: iteration convergence tolerance is 0.05 MW (0.05 MVar), the upper limit of power generation at

node 2 is 32 MW, active power flows in all lines should be within transfer capability. The difference between the two vectors is measured by the Euclidean distance and the Chebyshev distance. As a result of applying this method, the load flow solution with the maximum possible loads at nodes 3 and 4 is calculated. Figure 2 shows the values of active power injection which are the results of the calculation of three load flow solutions: the obtained loads at nodes 3 and 4 are equal to the forecast loads (Steady state); the obtained loads at nodes 3 and 4 are the maximum loads according to Euclidean distance (SSeuclidean) and Chebyshev distance (SSchebyshev), respectively.

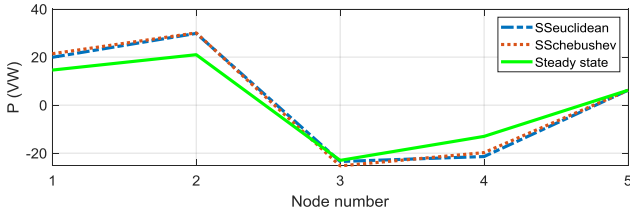


Fig. 2. The active power injections

5.3 Nonlinear optimization applying

Solving the problem which applies the nonlinear optimization, developed in Matlab, can be divided into several steps.

Step 1. Calculate coefficients (k_{i-j}) according to (21).

Step 2. Describe the objective function and constraints in the equivalent forms which are suitable for the programs, developed in Matlab.

The objective function

$$(\bar{P}_2 - P_2(z_2)) + (\bar{P}_3 - P_3(z_3)) + (\bar{P}_4 - P_4(z_4)) \rightarrow \max \quad (29)$$

in an equivalent form is written as follows:

$$-z_2 F_2 - z_3 F_3 - z_4 F_4 \rightarrow \max. \quad (30)$$

Similar transformations are performed for constraints.

EPS active power balance:

$$P_1 + P_2(z_2) + P_5 - P_3(z_3) - P_4(z_4) = 0 \quad (31)$$

is:

$$F_2 Z_2 + F_3 Z_3 + F_4 Z_4 = P_3^{max} + P_4^{max} - P_2^{max} - P_1 - P_5 \quad (32)$$

Active power balance at nodes 3,4

$$\Delta P_3 + \Delta P_{3-1} + \Delta P_{3-4} = 0, \quad (33)$$

$$\Delta P_4 + \Delta P_{4-2} + \Delta P_{4-3} = 0, \quad (34)$$

are transformed into the equations:

$$F_2 Z_2 (k_{3-1} + k_{3-4}) + F_3 Z_3 = (P_3 - P_3^{max}) + (P_2 - P_2^{max})(k_{3-1} + k_{3-4}) + b. \quad (35)$$

$$F_2 Z_2 (k_{4-2} + k_{3-4}) + F_4 Z_4 = (P_4 - P_4^{max}) + (P_2 - P_2^{max})(k_{4-2} + k_{3-4}) + b. \quad (36)$$

Transmission constraints

$$P_{i-j} + k_{i-j} \Delta P_2(z_2) < P_{i-j}^{max} \quad (37)$$

are transformed into the following inequalities:

$$F_2 k_{1-2} z_2 < P_{1-2}^{max} - P_{1-2} - k_{1-2} (P_2 - P_2^{max}), \quad (38)$$

$$F_2 k_{1-3} z_2 < P_{1-3}^{max} - P_{1-3} - k_{1-3} (P_2 - P_2^{max}), \quad (39)$$

$$F_2 k_{1-5} z_2 < P_{1-5}^{max} - P_{1-5} - k_{1-5} (P_2 - P_2^{max}), \quad (40)$$

$$F_2 k_{2-4} z_2 < P_{2-4}^{max} - P_{2-4} - k_{2-4} (P_2 - P_2^{max}), \quad (41)$$

$$F_2 k_{3-4} z_2 < P_{3-4}^{max} - P_{3-4} - k_{3-4} (P_2 - P_2^{max}). \quad (42)$$

Optimization parameters constraints are:

$$0 \leq z_2 \leq 1, \quad (43)$$

$$0 \leq z_3 \leq 1, \quad (44)$$

$$0 \leq z_4 \leq 1. \quad (45)$$

The objective function in the Matlab codes is:

[z,fval]=fmincon(@funn,z0,ineq_l,ineq_r,A,B,zmin,zmax);

Function f= funn; $f = -F_2 z_2 - F_3 z_3 - F_4 z_4$; initial approximation of optimization parameters is :z0 = [0 0 0]. A compact matrix formulation is used for representing constraints.

Equality constraints (A, B) are:

| A | B |
|---------------------------|--|
| $F_2 k_{1-5}$ | $0 \quad 0 \quad (P_2 - P_2^{max}) k_{1-5}$ |
| $F_2 (k_{1-3} + k_{3-4})$ | $F_3 \quad 0 \quad \bar{P}_3 - P_3^{max} + (P_2 - P_2^{max})(\sum_{j=1}^G k_{ij})$ |
| $F_2 (k_{4-2} + k_{3-4})$ | $0 \quad F_4 \quad \bar{P}_4 - P_4^{max} + (P_2 - P_2^{max})(\sum_{j=1}^G k_{ij})$ |
| F_2 | $F_3 \quad F_4 \quad \sum_1^{u-un} P_j^{max} + \sum_1^{un} P_i$ |

Inequality constraints (L, R,) are:

| L | R |
|---------------|---|
| $F_2 k_{1-2}$ | $0 \quad 0 \quad P_{1-2}^{max} - P_{1-2} - k_{1-2} (P_2 - P_2^{max})$ |
| $F_2 k_{1-3}$ | $0 \quad 0 \quad P_{1-3}^{max} - P_{1-3} - k_{1-3} (P_2 - P_2^{max})$ |
| $F_2 k_{1-5}$ | $0 \quad 0 \quad P_{1-5}^{max} - P_{1-5} - k_{1-5} (P_2 - P_2^{max})$ |
| $F_2 k_{2-4}$ | $0 \quad 0 \quad P_{2-4}^{max} - P_{2-4} - k_{2-4} (P_2 - P_2^{max})$ |
| $F_2 k_{3-4}$ | $0 \quad 0 \quad P_{3-4}^{max} - P_{3-4} - k_{3-4} (P_2 - P_2^{max})$ |

Step 3. Perform optimization. The result is a vector of optimization parameters (z) with given constraints.

Step 4. Interpret the results. Calculate the load values at nodes 3, 4, and the generation at node 2, according to formula (6). Figure 3 shows the active power values at nodes 2, 3, 4 before (SS), and after (SS max) optimization.

Step 5. Check if the obtained state variables meet the given limits. If the result is negative, invalid state variables are assumed to be corrected.

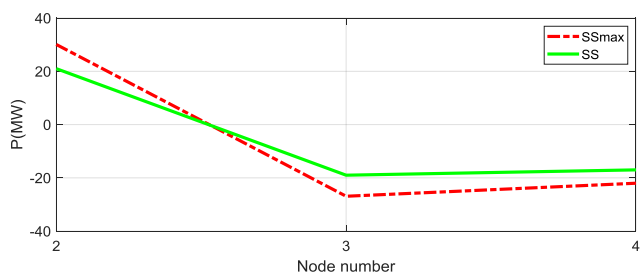


Fig. 3 Active power at nodes 2, 3, 4 before (SS) and after (SS max) optimization

5.4 Analysis of results

The results obtained by the two methods are summarized in Table 2. The last line shows the flexibility calculated by (28) using different methods. Table 3 presents the values calculated by the formula:

$$\Delta F_{dist} = |F_{nonl} - F_{dist}| \quad , \quad (46)$$

where F_{dist} is the EPS flexibility (F_s Table 2) calculated by the brute-force method using Euclidean distance (8.8MW) or Chebyshev distance (9.3MW), F_{nonl} is EPS flexibility (12.9 MW, Table 2) calculated by nonlinear optimization.

Table 2. Result of calculations (MW).

| Number of nodes | Initial data | | Calculated data | | |
|-----------------|------------------------|-----------------|--------------------------|--------------|------------------------|
| | Forecast p_i^{forec} | Max p_i^{max} | Brute force p_i^{calc} | | Nonlinear p_i^{calc} |
| | | | euclid | cheb | |
| 1 | 13.9 | | | | |
| 2 | 21 | 30 | 31 | | 30 |
| 3 | -23 | 37 | -23.4 | -25.5 | -26.9 |
| 4 | -13 | 25 | -21.4 | -19.8 | -22 |
| 5 | 6.2 | 6.2 | 6.2 | | |
| F_s | | | 8.8 | 9.3 | 12.9 |

Table 3. Absolute difference between two values of flexibility.

| Difference | $\Delta F_{dist(t)}$ (MW) | | $\sum \Delta F_{dist}$ (MW) |
|-----------------------|---------------------------|-----|-----------------------------|
| | Number of nodes | | |
| | 3 | 4 | |
| $F_{nonl} - F_{eucl}$ | 3.5 | 0.6 | 4.1 |
| $F_{nonl} - F_{cheb}$ | 1.4 | 2.2 | 3.6 |

The active powers, which are the result of three load flow solution problems and the result of the nonlinear optimization, are shown in Figure 4.

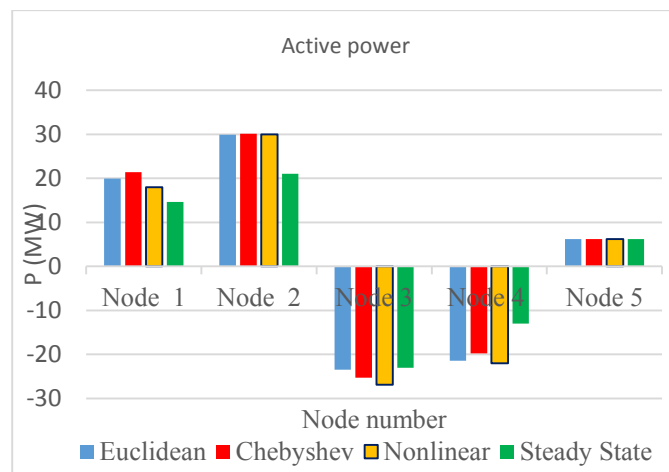


Fig. 4. Diagram of active power

Analysis of the results presented in Table 2 and Figure 4 shows that

- Comparison of the given maximum loads (italics) with the maximum possible loads calculated by different methods (bold) shows that the loads calculated using nonlinear optimization are closer to the given maximum loads;
- The flexibility of the considered EPS is 12.9 MW (formula 28) according to nonlinear optimization, 8.8 MW and 9.3 MW according to the brute-force method when using the Euclidean distance and Chebyshev distance, respectively, as a metric.

Table 3 shows that the nonlinear optimization results are closer to the results obtained by the brute-force method in the case of the Chebyshev distance used as a metric (4.1 > 3.6).

Findings have revealed that the calculation of load flow solution in the case where the result of nonlinear optimization is used as the initial data, requires that reactive power be added at node 3 for all variables to be within given limits.

6 Conclusion

The paper describes the methods for determining the flexibility of the electric power system: the brute-force and nonlinear optimization. In the brute-force method, the Euclidean distance and the Chebyshev distance are used as a metric for comparing the two vectors. For nonlinear optimization, a function developed in Matlab is used.

The nonlinear optimization results are analyzed. The analysis shows that the loads calculated using nonlinear optimization are closer to the given maximum loads. The study indicates that to ensure the reactive power balance in the EPS at the obtained load values, it is necessary to increase the reactive power at node 3.

A comparative analysis of the results has shown that the reference for the verification of the nonlinear optimization output should be the results calculated by the brute-force algorithm based on the Chebyshev distance as a metric.

An algorithm was developed to create an archive of loads required for the brute-force method.

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